

Limits of network nonlocality probed by time-like separated observers

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In an entanglement swapping scenario, if two sources sharing entangled states between three parties are independent, local correlations lead to a different kind of inequalities than the standard Bell inequalities, known as network local models. A highly demanding task is to find out a way to involve many players nontrivially in a quantum network since measurements, in general, disturb the system. To this end, we consider here a novel way of sharing network nonlocality when two observers initially share close to a maximally entangled state. We report that by employing unsharp measurements performed by one of the observers, six pairs can sequentially demonstrate the violation of bilocal correlations while a maximum of two pairs of observers can exhibit bi-nonlocality when both the observers perform unsharp measurements. We also find the critical noise involved in unsharp measurements in each round to illustrate the bi-nonlocality for a fixed shared entangled state as a resource. We also establish a connection between entanglement content of the shared state, quantified via von-Neumann entropy of the local density matrix for pure states and entanglement of formation for Werner states, and the maximum number of rounds showing violation of bilocal correlations. By reducing entanglement content in the elements of the joint measurement by the third party, we observe that the maximum number reduces to two sequential sharing of bi-nonlocality even for the maximally entangled state when the settings at each side are taken to be three and fixed.

I. INTRODUCTION

Quantum resource states are shown to enhance capacities in transmitting both classical and quantum information over classically known protocols which were later implemented successfully between a single sender and receiver using different physical substrates [1–6]. However, quantum technological developments also require generalization and realization of these protocols in a multipartite domain involving several parties situated in distant locations, thereby building a quantum communication network [7]. One of the prominent designs in this direction is the proposal of quantum repeaters, a combination of entanglement distillation and swapping, [8–11] by which entangled states are shared between observers separated by long-distance even in presence of noise. A crucial step here is to verify the resource content in the created states. For shared entangled states, several identification schemes exist which include testing Bell inequalities [12, 13], entanglement witnesses [14], steering inequality [15].

Apart from entanglement detection, studies of Bell's theorem plays an important role in understanding quantum theory [12, 13]. It was shown that bipartite entangled pure states always violate some Bell inequalities [16]. Over the years, Bell inequalities have also been generalized in multipartite domains and hence become crucial to establish nonlocality in networks [13, 17–19]. It has been realized that if one considers that the sources which share entangled states are independent, a distinct kind of local realistic models can be constructed which are different from the standard Bell inequalities – a violation of these inequalities confirms the nonlocality in networks [20–24]. The simplest network is called the bilocal scenario involving two independent sources which share two entangled states between three parties having three inputs and outputs – a violation of the inequality that confirms the impossibility of local models was introduced by Branciard-

Rosset-Gisin-Pironio, referred to as BRGP inequality. In these scenarios, several works have been carried out both in the chain of arbitrary length and in star networks for which different kinds of inequalities based on local models can be derived. Moreover, unlike the paradigms of standard Bell inequalities, independent resource consideration leads to much more involved structures in the set of local correlations which include nonconvexity of the set.

On the other hand, projective or sharp measurements can reveal nonlocal correlations present in the states by collecting the statistics required for Bell inequalities although they can destroy the shared entanglement. On the contrary, weak or unsharp measurements can serve both purposes by providing a trade off relation between information gain and disturbance due to measurement [25]. In recent times, the generalized measurements are shown to be important tools in various quantum information tasks like state discrimination [26], state tomography [27], violation of Bell inequalities [28], randomness generation [29], detection of entanglement [30], creating multipartite entangled states [31]. At the same time, it was also found that a shared entangled state can be detected sequentially by the violation of Bell and steering inequalities, device-dependent as well as -independent entanglement witnesses by a single observer or by both the observers where observers perform unsharp measurements [32–41]. Upto now, all the sequential scenarios considered assume that there is a single source which produces the shared state initially.

In the present work, we go beyond this picture (cf. [42, 43]). In particular, we consider two independent sources which produce two noisy nonmaximally entangled states. In this situation, after joint measurement by the middle party, the other two observers' aim is to check nonlocal correlations and in the sharing scenario, the task is done by unsharp measurements (see Fig. 1). In this paper, we consider two sequential scenarios – (1) depending on the unsharp measurement per-

formed by one of the observers, they exhibit nonlocal correlations sequentially by obtaining violations of bilocal inequalities which we call unidirectional sharing of bi-nonlocality; (2) both the observers find the critical unsharp parameters in each round to manifest nonlocal correlations which we refer to as bidirectional sharing of bi-nonlocality. In both scenarios, we establish a connection between the entanglement content of the shared state and the maximum number of cycles in which they are capable to demonstrate bi-nonlocal correlations. We observe that in the unidirectional case, unlike standard Bell inequalities, the violation of bilocal inequalities can be observed with a maximum of *six* rounds when the shared state is close to maximally entangled states, both for pure and noisy states. The maximum number reduces to two when both the observers wish to demonstrate the bi-nonlocality.

In an entanglement swapping protocol, the middle party has to carry out a joint entangling measurement on his/her parts. All the above results are obtained when Bell-basis measurements are performed by the middle party. In contrast, if the middle party performs a more general joint measurement, known as elegant joint measurement, another kind of inequality emerges to detect nonlocal correlations having three input settings [44]. We show that in this scenario, the maximum number reduces to two even for a unidirectional case with the noisy entangled state having high entanglement as an initial resource while we obtain that the sequential sharing is not possible when both the observers perform unsharp measurements, thereby reaching to a no-go theorem. Note that the results are true when each element in the elegant joint measurement basis contain a minimum amount of entanglement.

The paper is organized in the following way. In Sec. II, we introduce the bilocal inequalities both for Bell-basis and elegant joint measurements. We first present the recursion relation of different rounds involved in sequential sharing (SubSec. III A) and two sequential scenarios, unidirectional (SubSec. III B) and bidirectional ones (SubSec. III C) when the shared state is maximally entangled pure states. We then consider the sequential scenario with noisy nonmaximally entangled states in Sec. IV, thereby establishing a relation between entanglement of the shared state and the maximal number of observers exhibiting network nonlocality. Going beyond Bell-basis measurement, and considering elegant joint measurement, the sharing scenario changes drastically which will be discussed in Sec. V. We finally conclude in Sec. VI.

II. NETWORK INEQUALITIES WITH DIFFERENT JOINT MEASUREMENTS

Let us briefly describe the network nonlocality, which is different from the standard Bell inequality. In an entanglement swapping scenario [9, 10], we assume that a single source creates two copies of a bipartite state, ρ , which are shared between Alice-Bob (AB) and Bob-Charu (BC) pairs. After Bob's joint measurement on his parts, Alice and Charu can share an entangled state, ρ' whose entanglement content depends on the initial pairs and joint measurements by B . Note that if the initial states are maximally entangled, the

Bell-basis measurement at B 's node projects the state between A and C into maximally entangled. To detect entanglement between A and C , several methods can be employed which include entanglement witness [14], standard Bell inequalities [12, 13], steering inequality to name a few.

Instead of a single source, we now assume that there are two independent sources S_1 and S_2 , which emit two states characterized by hidden variables λ_1 and λ_2 respectively and states corresponding to λ_1 (λ_2) is shared by AB (BC) as shown in Fig. 1. Here A and C have measurement settings labeled by x and z with outcomes a and c , respectively while B has a fixed measurement setting. In this situation, a new paradigm emerges known as bilocal scenario [20–23]. In this paper, we consider two kinds of measurements performed by B – 1. *Bell-basis measurement* (BSM), and 2. *elegant joint measurement* (EJM). We will discuss about these bases in later part of this section. It has been established that for BSM, any bilocal model has to satisfy BRGP inequality [21] while different bilocal inequality is derived in case of EJM [44]. We will now briefly discuss both of them.

A. Local models based on BSM

Let A and C have binary inputs and outputs, i.e., $x, z, a, c \in \{0, 1\}$. However, B has two bits of output, $\mathbf{b} = b^0 b^1 = 00, 01, 10, 11$ corresponding to the Bell-basis measurement, $\{|\phi^\pm\rangle = \frac{1}{\sqrt{2}}(|00\rangle \pm |11\rangle), |\psi^\pm\rangle = \frac{1}{\sqrt{2}}(|01\rangle \pm |10\rangle)\}$ and the projectors for BSM are denoted by $\Pi_{b^0 b^1}$. From the conditional probability, $P^{14}(a, b^0 b^1, c|x, z)$ obtained by three parties A, B and C , after measurements being performed on their parts, let us define the tripartite correlation, given by

$$\langle A_x B^y C_z \rangle = \sum_{a, b^0 b^1, c} (-1)^{a+b^y+c} P^{14}(a, b^0 b^1, c|x, z), \quad (1)$$

where $y \in \{0, 1\}$. Taking linear combinations of the above correlations, we construct two quantities, represented as

$$I^{14} = \frac{1}{4} \sum_{x, z} \langle A_x B^0 C_z \rangle, \quad (2)$$

$$J^{14} = \frac{1}{4} \sum_{x, z} (-1)^{x+z} \langle A_x B^1 C_z \rangle. \quad (3)$$

It was shown [20, 21] that any bilocal model based on two independent sources would satisfy the inequality, given by

$$\mathcal{B} := \sqrt{|I^{14}|} + \sqrt{|J^{14}|} \leq 1. \quad (4)$$

We will refer to the left hand side as BRGP function or expression.

B. Different bilocal scenario with EJM

Instead of Bell-basis measurements by B , we now consider a scenario in which B performs a joint entangling measurement given by $\{|\Psi_b^\theta\rangle\}_{b=1}^4$, parametrized by $\theta \in \{0, \frac{\pi}{2}\}$ with

all elements of the basis being equally entangled [44]. To construct the basis with this property, let us first write the pure states in cylindrical coordinates, representing the four vertices in a regular tetrahedron inside the Bloch sphere as

$$|\pm\vec{m}_b\rangle = \sqrt{\frac{1 \pm r_b}{2}} e^{-i\frac{\phi_b}{2}} |0\rangle \pm \sqrt{\frac{1 \mp r_b}{2}} e^{i\frac{\phi_b}{2}} |1\rangle. \quad (5)$$

With the help of them, EJM basis reads as

$$|\Psi_b^\theta\rangle = \frac{\sqrt{3} + e^{i\theta}}{2\sqrt{2}} |\vec{m}_b, -\vec{m}_b\rangle + \frac{\sqrt{3} - e^{i\theta}}{2\sqrt{2}} |-\vec{m}_b, \vec{m}_b\rangle. \quad (6)$$

Notice that by varying θ from 0 to $\pi/2$, one can reach from EJM to BSM (upto some local unitaries).

Unlike BSM, A and C can choose to perform measurement out of three possible settings for each of them, i.e., $x, z \in \{0, 1, 2\}$, with binary outcomes $a, c \in \{0, 1\}$. The four possible output of B , representing the four vertices of the tetrahedron, \vec{m}_b , given by $\vec{m}_1 = (1, 1, 1)$, $\vec{m}_2 = (1, -1, -1)$, $\vec{m}_3 = (-1, 1, -1)$ and $\vec{m}_4 = (-1, -1, 1)$ can be labelled as the three-vector $\mathbf{b} = (b^1, b^2, b^3)$. In joint measurement by B , the bilocal inequality reads as

$$\mathcal{B}_E := \frac{1}{3} \left(\sum_{y=z} \langle B^y C_z \rangle - \sum_{x=y} \langle A_x B^y \rangle \right) - \sum_{x \neq y \neq z \neq x} \langle A_x B^y C_z \rangle \leq 3 + 5Z. \quad (7)$$

where $Z = \max \{ |\langle A_x \rangle|, |\langle A_x B^y \rangle|, \dots, |\langle A_x B^y C_z \rangle| \}$ is the maximum of the absolute values of marginal and full correlators, which do not appear in \mathcal{B}_E . As obtained in case of BSM in Eq. (1), we can also write the above correlators in terms of the conditional probabilities emerged from experiments. For example, we have

$$\langle A_x B^y C_z \rangle = \sum_{a, b^1, b^2, b^3, c} b^y (-1)^{a+c} p(a, \mathbf{b}, c | x, z). \quad (8)$$

The left hand side of (7) can be called Tavakoli-Gisin-Branciard (TGB) function.

III. SEQUENTIAL DETECTION OF BI-NONLOCALITY

Let us now set the framework of sharing bipartite quantum states sequentially. In this work, we consider following two scenarios (see Fig. 1 for schematics) –

A. Unidirectional. One of the parties performs unsharp measurements and the other spatially separated observer does projective measurement after B 's joint measurement in an entanglement swapping experiment.

B. Bidirectional. Both the parties, i.e., A s and C s perform weak measurements, thereby disturbing the state minimally which occurs after the join measurement is completed at B 's end.

Before presenting the results, let us discuss the general protocol that will be followed to sequentially share and test quantum network nonlocality.

A. General protocol for sharing bi-nonlocality in network

After Bob's Bell-basis measurement, Alices (Charus) denoted as A^1, A^2, \dots, A^m (C^1, C^2, \dots, C^n) perform weak measurements and send their parts of the qubits to the next Alice (Charu) which can capture the competition between information gain and disturbance due to measurement. The measurement choices of Alices (Charus) as inputs can be denoted as x_1, x_2, \dots, x_m (z_1, z_2, \dots, z_n) and the corresponding measurement outcomes can be denoted as a_1, a_2, \dots, a_m (c_1, c_2, \dots, c_n). As mentioned before, $\{x_m\}, \{z_n\}, \{a_m\}, \{c_n\} \in \{0, 1\}$. Suppose A^m and C^n choose their measurement directions with angle ϕ_m and θ_n respectively in the $x - z$ plane, given by

$$\begin{aligned} A_{x_m}^m &= \cos \phi_m \sigma_z - (-1)^{x_m} \sin \phi_m \sigma_x, \\ C_{z_n}^n &= \cos \theta_n \sigma_z + (-1)^{z_n} \sin \theta_n \sigma_x. \end{aligned} \quad (9)$$

The projectors of Alices and Charus corresponding to their inputs and outputs are respectively written as

$$\begin{aligned} \Pi_{x_m}^{a_m} &= \frac{(\mathbb{I} + (-1)^{a_m} A_{x_m}^{a_m})}{2} \otimes \mathbb{I}, \\ \Pi_{z_n}^{c_n} &= \mathbb{I} \otimes \frac{(\mathbb{I} + (-1)^{c_n} C_{z_n}^{c_n})}{2}. \end{aligned} \quad (10)$$

Using these tools and concepts of unsharp measurements discussed in Appendix. A, we can put forward some simple steps to find the joint probability distribution $P(a_m, b^0 b^1, c_n | x_m, z_n)$, which will be required to check BRGP inequality (bilocal model with EJM) of the output states in each round. Let us enumerate each step of the protocol by Alices and Charus in details.

- Initial state shared by Alice-Bob-Charu is denoted by $\rho = \rho_{A^1 B} \otimes \rho_{B C^1}$. After Bob performs the Bell-basis measurement, the shared state between Alice and Charu depending on Bob's outcome can be written as

$$\rho_{A^1 C^1}^{b^0 b^1} = \text{tr}_B [(\mathbb{I} \otimes \Pi_{b^0 b^1} \otimes \mathbb{I}) \rho (\mathbb{I} \otimes \Pi_{b^0 b^1} \otimes \mathbb{I})^\dagger]. \quad (11)$$

- Each Alice (till $(m-1)$ Alice, i.e., A^{m-1}) performs unsharp measurement according to the choice of the string x_1, x_2, \dots, x_{m-1} and the corresponding quality factors of the weak measurements, F_1, F_2, \dots, F_{m-1} [32]. For a fixed round, say, k , after A^k 's measurement, the part of the state is sent to the next Alice, i.e., A^{k+1} without communicating the outcome. Hence the final transformed state between A^m and C^1 after $m-1$ rounds of measurement by previous Alices, depending on all the previous Alice's measurement choices, can be written as

$$\begin{aligned} \rho_{A^m C^1}^{b^0 b^1 | x_1, x_2, \dots, x_{m-1}} &= \\ \mathcal{W}_{x_{m-1}} (\mathcal{W}_{x_{m-2}} (\dots \mathcal{W}_{x_1} (\rho_{A^1 C^1}^{b^0 b^1}) \dots)), \end{aligned} \quad (12)$$

where the map \mathcal{W}_{x_i} is defined as

$$\begin{aligned} \mathcal{W}_{x_i}(\rho) &= \\ = F_i \rho + (1 - F_i) (\Pi_{x_i}^0 \rho (\Pi_{x_i}^0)^\dagger + \Pi_{x_i}^1 \rho (\Pi_{x_i}^1)^\dagger). \end{aligned} \quad (13)$$

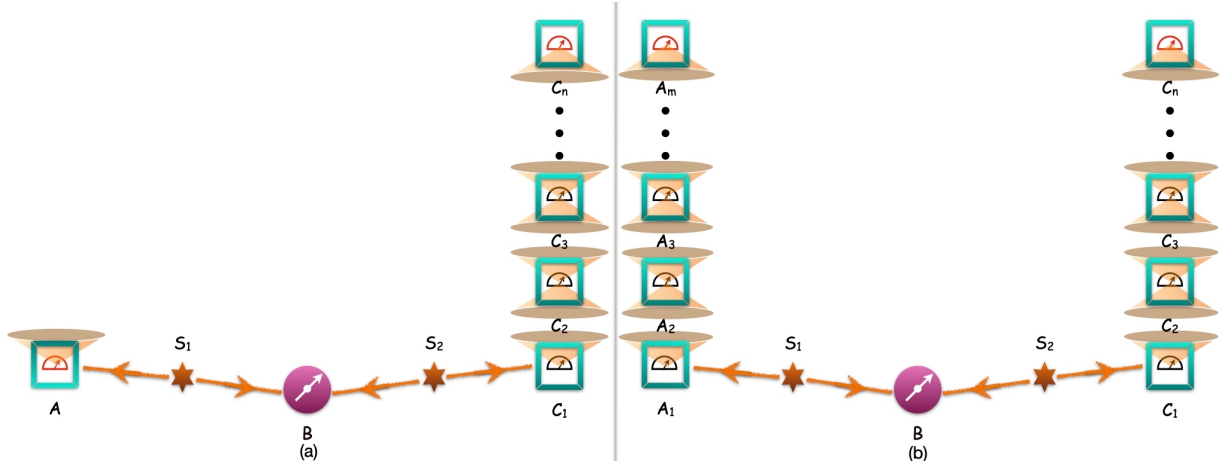


FIG. 1. (Color online.) Schematic diagram for sharing of quantum states sequentially. (a) Unidirectional where one of the observers performs weak measurement. (b) Bidirectional sharing of states in which both the observers perform weak measurements. In both the scenarios, nonlocal nature of states in each round is confirmed from the violation of network nonlocality (bilocal inequality).

- In a similar fashion, each Charu (till $(n-1)$ Charu, i.e., C^{n-1}) performs unsharp measurement according to the choice of the string z_1, z_2, \dots, z_{n-1} and quality factors $F'_1, F'_2, \dots, F'_{m-1}$. The resulting state between A^m and C^n in this case reads as

$$\rho_{A^m C^n}^{b^0 b^1 | x_1, x_2, \dots, x_{m-1}, z_1, z_2, \dots, z_{n-1}} = \quad (14)$$

$$\mathcal{W}_{z_{n-1}}(\mathcal{W}_{z_{n-2}}(\dots \mathcal{W}_{z_1}(\rho_{A^m C^1}^{b^0 b^1 | x_1, x_2, \dots, x_{m-1}}) \dots)),$$

where

$$\mathcal{W}_{z_i}(\rho) = F'_i \rho + (1 - F'_i)(\Pi_{z_i}^0 \rho (\Pi_{z_i}^0)^\dagger + \Pi_{z_i}^1 \rho (\Pi_{z_i}^1)^\dagger). \quad (15)$$

- In the last step, A^m and C^n perform unsharp measurement according to the measurement choice x_m , and z_n with outcome a_m , and c_n respectively. The post measurement state becomes

$$\rho_{z_n}^{a_m, b^0 b^1, c_n | x_1, x_2, \dots, x_m, z_1, z_2, \dots, z_n} \quad (16)$$

$$= \mathcal{W}_{z_n}^{c_n}(\mathcal{W}_{x_m}^{a_m}(\rho_{A^m C^n}^{b^0 b^1 | x_1, x_2, \dots, x_{m-1}, z_1, z_2, \dots, z_{n-1}})).$$

where the corresponding operators, $\mathcal{W}_{x_m}^{a_m}$ and $\mathcal{W}_{z_n}^{c_n}$ can be represented as

$$\begin{aligned} \mathcal{W}_{x_m}^{a_m}(\rho) & \quad (17) \\ &= \frac{F_m}{2} \rho + \frac{(1 + (-1)^{a_m} G_m - F_m)}{2} (\Pi_{x_m}^0 \rho (\Pi_{x_m}^0)^\dagger) \\ &+ \frac{(1 - (-1)^{a_m} G_m - F_m)}{2} (\Pi_{x_m}^1 \rho (\Pi_{x_m}^1)^\dagger). \end{aligned}$$

We are now ready to compute the joint probability distribution, given by

$$\begin{aligned} P(a_m, b^0 b^1, c_n | x_1, x_2, \dots, x_m, z_1, z_2, \dots, z_n) & \quad (18) \\ &= \text{tr}(\rho_{z_n}^{a_m, b^0 b^1, c_n | x_1, x_2, \dots, x_m, z_1, z_2, \dots, z_n}). \end{aligned}$$

- We also require to consider the previous individual probabilities by Alices and Charus measurement choices and by performing average over all such measurement choices, we obtain

$$P^{14}(a_m, b^0 b^1, c_n | x_m, z_n) \quad (19)$$

$$= \sum_{x_1, \dots, x_{m-1}, z_1, \dots, z_{n-1}=0}^1 P(x_1) \dots P(x_{m-1}) P(z_1) \dots$$

$$P(z_{n-1}) \times P(a_m, b^0 b^1, c_n | x_1, x_2, \dots, x_m, z_1, z_2, \dots, z_n)$$

$$= \sum_{x_1, \dots, x_{m-1}, z_1, \dots, z_{n-1}=0}^1 \frac{1}{2^{n+m-2}} \times$$

$$P(a_m, b^0 b^1, c_n | x_1, x_2, \dots, x_m, z_1, z_2, \dots, z_n)$$

- After finding these joint correlations between A^m, B , and C^n , it is straightforward to find the conditions for which these correlations violate bilocal models using BRGP inequality. Specifically, the quantities required for BRGP inequality take the form as

$$\langle A_{x_m}^m B^y C_{z_n}^n \rangle \quad (20)$$

$$= \sum_{a_m, b^0 b^1, c_n} (-1)^{a_m + b^y + c_n} P^{14}(a_m, b^0 b^1, c_n | x_m, z_n),$$

$$I_{m,n}^{14} = \frac{1}{4} \sum_{x_m, z_n} \langle A_{x_m}^m B^y C_{z_n}^n \rangle,$$

$$J_{m,n}^{14} = \frac{1}{4} \sum_{x_m, z_n} (-1)^{x_m + z_n} \langle A_{x_m}^m B^y C_{z_n}^n \rangle,$$

and finally we obtain the condition on (G, F) -pair such that

$$\mathcal{B}(A^m, B, C^n) := \sqrt{|I_{m,n}^{14}|} + \sqrt{|J_{m,n}^{14}|} > 1. \quad (21)$$

In case of unidirectional sharing, our aim is to find maximum m or n by performing projective measurement by the other party, i.e., by fixing the other index to be 1. On the other hand, the maximum pair of (m, n) will be found for the bidirectional situation. Notice also that instead of Bell-basis measurement, if B performs EJM, we can also compute the corresponding network inequality in each round by slightly modifying all the derivations obtained above.

B. Unidirectional sharing of bi-nonlocality

In the case of unidirectional sharing, our motivation is to find the criteria under which all temporally separated Charus share a bi-nonlocal correlation with a single Alice, thereby violating the BRGP inequality. Without loss of generality, we can take $m = 1$ and hence we assume that A^1 does not perform weak measurement, i.e., she performs a sharp measurement with $G_1 = 1$.

Let us first illustrate the situation when the shared states between (A^1, B) and (B, C^1) pairs are maximally entangled and B performs the Bell-basis measurement. After some manipulations, the general form of BRGP inequality between A^1, B and C^n can be written as

$$\mathcal{B}(A^1, B, C^n) = \sqrt{|I_{1,n}^{14}|} + \sqrt{|J_{1,n}^{14}|}, \quad (22)$$

with

$$I_{1,n}^{14} = G'_n \cos \theta_n \cos \phi \sum_{\{l_i\}=0}^1 \prod_{i=1}^{n-1} (1 + (-1)^{l_i} F'_i) (\cos 2\theta_i)^{l_i},$$

$$J_{1,n}^{14} = G'_n \sin \theta_n \sin \phi \sum_{\{l_i\}=0}^1 \prod_{i=1}^{n-1} (1 + (-1)^{l_i} F'_i) (\cos 2\theta_i)^{l_i} (-1)^{l_i}$$

(23)

The expressions for a few rounds are mentioned in Appendix A 1 which can clearly give us the idea to obtain optimal rounds for sharing.

1. Optimal strategy to share bi-nonlocality

After obtaining the BRGP expression in Eq. (23), for a fixed round, say, k , we compute the minimum G'_k value for which $\mathcal{B}(A^1, B, C^k)$, a function of G'_k, θ_k and ϕ , just starts violating the BRGP inequality. It can be easily confirmed by considering Eq. (A4) that minimum G'_k is obtained for any round, k , when $\theta_k = \phi = \pi/4$. For example, $\mathcal{B}(A^1, B, C^1)=1$ leads to condition for critical G_1^{cr} as

$$\sqrt{G_1^{cr}} = \frac{1}{\sqrt{|\cos \theta_1 \cos \phi|} + \sqrt{|\sin \theta_1 \sin \phi|}} \quad (24)$$

Now it is obvious that $\min(\sqrt{G_1^{cr}}) = \frac{1}{\sqrt{2}}$ at $\theta_1 = \phi = \frac{\pi}{4}$. Putting these values of G_1^{cr}, θ_1, ϕ in $\mathcal{B}(A^1, B, C^2)$ and demanding it to be unity, we can similarly show $\sqrt{G_2^{cr}}$ is minimum with $\theta_2 = \frac{\pi}{4}$. Same arguments apply for all rounds

and finally we can specify the optimal strategy of measurement at each round such that the state is minimally disturbed or probed to show bi-nonlocality with all $G'_i > G_i^{cr}$. Under this conditions, the general form of $\mathcal{B}(A^1, B, C^n)$ simplifies as

$$\mathcal{B}(A^1, B, C^n) = 2 \sqrt{\frac{1}{2^n} \prod_{i=1}^{n-1} (1 + F'_i) G'_n}. \quad (25)$$

■ Proposition I. *In the unidirectional sharing of bi-nonlocality, when two independent sources produce two copies of maximally entangled states, a single Alice can simultaneously violate BRGP inequality with a maximum of six Charus provided Bell-basis measurement is performed by the middle party (Bob).*

Proof. Let the maximum number of Charu showing bi-nonlocality with Alice be n . To find the critical values of $G_i^{cr}, \forall i = 1, \dots, n$, we need the solutions of $\mathcal{B}(A^1, B, C^n) = 1, \forall i = 1, \dots, n$ which lead to

$$G_1^{cr} = \frac{1}{2}, \quad (26)$$

$$G_{i+1}^{cr} = \frac{2G_i^{cr}}{1 + F_i^{cr}}. \quad (27)$$

Simplifying this, we get $G_2^{cr} = 0.536, G_3^{cr} = 0.581, G_4^{cr} = 0.64, G_5^{cr} = 0.725, G_6^{cr} = 0.859, G_7^{cr} = 1.135$. It immediately implies that only six Charus can satisfy Eq. (27) for $0 < G_i^{cr} \leq 1$ so that the resulting state violates the BRGP inequality. ■

In the succeeding sections, we will demonstrate that the BSM by Bob and entanglement content of the shared states are crucial to obtain the maximum number of rounds showing bi-nonlocality as depicted in Figs. 2 and 3.

Remark. The optimal measurement settings, i.e., the values of $\theta_k = \phi = \pi/4$ remain optimal also for the Werner state [45] $\forall k$ rounds (i.e., when the maximally entangled state is admixed with white noise).

C. Bidirectional sharing of bi-nonlocality: Advantage in asymmetry

As shown in the previous situation, the optimal choice of measurement direction in this case also turns out to be $\theta_i = \phi = \pi/4$. Using this, we can generalise the BRGP violation between m -th Alice and n -th Charu as

$$\mathcal{B}(A^m, B, C^n) = 2 \sqrt{\frac{1}{2^{n+m-1}} \prod_{i=1}^{m-1} (1 + F_i) \prod_{i=1}^{n-1} (1 + F'_i) G_m G'_n}. \quad (28)$$

In other words, m number of Alices and n number of Charus are said to be perfectly share bi-nonlocality bidirectionally if

$$\mathcal{B}(A^i, B, C^j) > 1 \quad \forall i = 1, \dots, m; j = 1, \dots, n. \quad (29)$$

Weak measurements with equal precision. In this scenario, if we take the precision of the measurement at Alice and Charu's end to be equal, i.e., $G'_n = G_n$, we have the following results.

■ **Proposition II.** *In the bidirectional sharing with equal precision in Alice and Charu's measurements in each round, a maximum number of Alice and Charu who can perfectly share bi-nonlocality sequentially is two when the shared state is maximally entangled and the Bell-basis measurement is performed.*

Proof. The proof is similar to Proposition I. With the equality condition in (29), we get the criteria as

$$G_1^{cr} = \frac{1}{\sqrt{2}}, \quad (30)$$

$$G_{i+1}^{cr} = \frac{2G_i^{cr}}{1 + F_i^{cr}}. \quad (31)$$

Calculating explicitly, we find $G_2^{cr} = 0.828$, $G_3^{cr} = 1.06$. Therefore, only two Alices can share perfect bi-nonlocality with two Charus having $0 < G_i^{cr} \leq 1$. ■

Weak measurements with unequal precision. Let us now take precision of unsharp measurements performed by Alices and Charus are unequal, i.e., $G_m \neq G_n$. It is interesting to check whether the situation is advantageous than the previous ones.

First, we find whether two Alices can share bi-nonlocality perfectly with more than two number of Charus. Without loss of generality, let us take sharpness parameter of second Alice to be 1. Now taking equality sign in (29), we obtain the following conditions, given by

$$G_2^{cr} = 1, \quad (32)$$

$$G_1^{tcr} G_1^{cr} = \frac{1}{2}, \quad (33)$$

$$G_{i+1}^{tcr} = \frac{2G_i^{tcr}}{1 + F_i^{tcr}}, \quad (34)$$

$$G_{i+1}^{cr} = \frac{2G_i^{cr}}{1 + F_i^{cr}}. \quad (35)$$

Using these relation, we immediately observe that $G_1^{tcr}, G_2^{tcr}, G_3^{tcr} < 1$ and $G_4^{tcr} > 1$. Thus in this asymmetric scenario, *two* Alices can show violation of BRGP inequality with *three* Charus. Similarly, for $m = 3(4)$, we can show the maximum number of Charus can be $n = 2(1)$. Hence, we prove that at most *two* Alices (Charus) can share bi-nonlocality with a maximum of *three* Charus (Alices) for a shared maximally entangled state and for BSM.

IV. DETECTING BI-NONLOCALITY SEQUENTIALLY WITH NOISY NONMAXIMALLY ENTANGLED STATES

Instead of sharing a two-qubit maximally entangled states between Alice, Bob and Charu, let us consider the situation when two independent sources can share noisy non-maximally entangled states written as $\rho = \rho_1^{AB} \otimes \rho_2^{BC}$ with $\rho_1^{AB} = v_1 |\psi_\alpha\rangle\langle\psi_\alpha| + \frac{1-v_1}{4}\mathbb{I}$, and $\rho_2^{BC} = v_2 |\psi_\beta\rangle\langle\psi_\beta| + \frac{1-v_2}{4}\mathbb{I}$,

having visibilities v_1 and v_2 respectively and $|\psi_\eta\rangle = \sqrt{\eta}|00\rangle + \sqrt{1-\eta}|11\rangle$ with $\eta = \alpha$ or β .

Following the similar prescription discussed in Sec. III A, we can generalize the BRGP function between Alice after m rounds and Charu after n rounds as

$$\mathcal{B}(A^m, B, C^n) = \sqrt{\frac{1}{2^{n+m-1}} \prod_{i=1}^{m-1} (1 + F_i) \prod_{i=1}^{n-1} (1 + F'_i) G_m G'_n} \times \sqrt{v_1 v_2} \times (1 + 2^4 \sqrt{\alpha(1-\alpha)\beta(1-\beta)}). \quad (36)$$

Using the above recursion relation, we can find the minimum disturbance value at each round so that the shared state can show network nonlocality in maximum rounds.

A. Bounds on sharing nonlocality between unidirectional time-like separated observers

NME as resource. Let us first manifest the maximum number of rounds for which bi-nonlocality can be shown when both AB and BC share identical copies of non-maximally entangled (NME) states, i.e., ρ_1^{AB} with $v_1 = 1$ (and similarly ρ_2^{BC} with $v_2 = 1$) and $\alpha = \beta$.

We want to examine the maximum number of sequential observers on Charu's (C^m) side can violate BRGP inequality with a single Alice (A^1) after B 's BSM. Using Eq. (37), we can show that a maximum of six Charus can sequentially demonstrate bi-nonlocality with Alice when the initial resource is close to a ME state.

The similar analysis also helps us to establish a connection between the entanglement content of the initial state, E_{in} , quantified by the von-Neumann entropy of the local density matrices [14] and the maximum number of rounds. Specifically, we find that with the decrease of entanglement in the initial resource states, the number of observers at one side (Charus) decreases as shown in Fig. 2. As shown in case of detection of entanglement sequentially via entanglement witness operators [34], we also observe that along with the maximally entangled state, there is also other non-maximally entangled states, having entanglement $E_{in} > 0.951$ which can exhibit violation of BRGP inequality upto six rounds with $\theta_i = \phi = \pi/4 \forall i$. Also, the sharing of nonlocality is possible (i.e., the maximum of two Charus can demonstrate the violation of BRGP inequality) for $E_{in} > 0.456$. It is to be noted that the hierarchy among NME states in this sharing scenario according to the violation of BRGP inequality is obtained with fixed settings, i.e., with $\theta_i = \phi = \pi/4$ ($i = 1, \dots, n$) (cf. [46]).

Noisy entangled states as resource. In the unidirectional domain, similar analysis can also be carried out by taking two identical copies of noisy entangled states as initial resources with $\alpha = \beta = 1/2$, i.e., the Werner states having $v_1 = v_2$ [45].

Interestingly, we report that there exists a critical noise value upto which Charu can show bi-nonlocality with a single Alice in maximum six rounds (see Fig. 3). In particular,

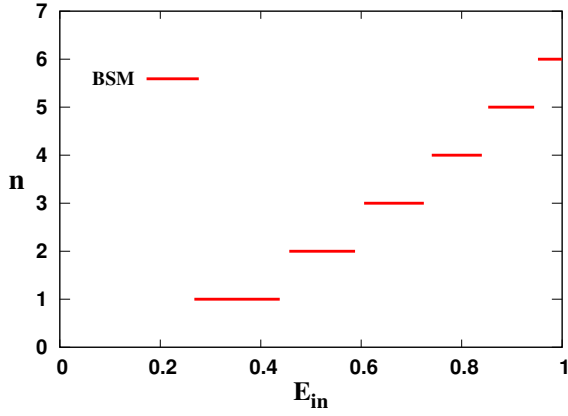


FIG. 2. (Color online.) **Maximum rounds vs. initial entanglement for NME states.** The abscissa, E_{in} denotes the entanglement of the initial resource states calculated in terms of von Neumann entropy of the local density matrices while the ordinate, n signifies the number of Charus who can show violation of BRGP inequality with $\theta_n = \phi = \pi/4$, thereby indicating the presence of noncoality in the shared pair. The finite length of the steps implies that the maximum number of Charus who can demonstrate nonlocality with Alice after Bob's Bell-basis measurement remains fixed for a finite range of initial entanglement. Both the axes are dimensionless.

if we calculate entanglement of formation EoF of the initial resource states [47], we find that when $EoF > 0.978$, the maximum rounds that Alice-Charu-duo can sequentially share states which violate BRGP inequality is six. On the other hand, when $EoF < 0.428$, Alice-Charu's state does not show violation even for a single round. The sequential protocol (i.e., minimum EoF above which two Charus can share bi-nonlocality with a single Alice) succeeds when $EoF > 0.591$.

V. SHARING BI-NONLOCALITY WITH EJM

Let us move to a scenario where Bob performs EJM given in Eq. (6) and the corresponding bilocal inequality also modifies as in Eq. (7). Initially, A^1B and BC^1 share the Werner states, $\rho_i = v_i |\psi^-\rangle\langle\psi^-| + \frac{1-v_i}{4}\mathbb{I}$ with visibility v_1 and v_2 respectively. Considering measurement settings for obtaining the violation of bilocal models as $\{\sigma_x, \sigma_y, \sigma_z\}$ for Alice and

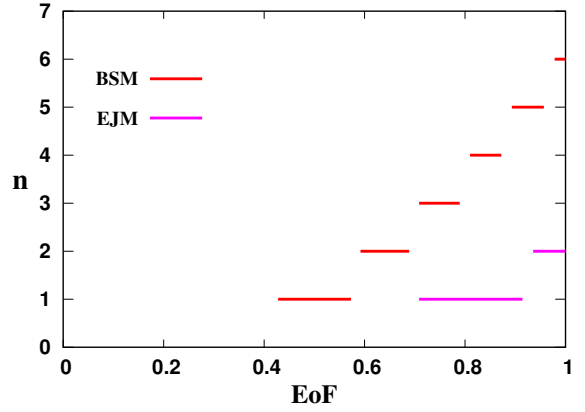


FIG. 3. (Color online.) **Maximum number of rounds against initial entanglement for Werner states as initials.** The maximum number of rounds, n by Charus (vertical axis) with respect to the initial entanglement quantified by entanglement of formation of the Werner state, EoF (horizontal axis). The implication is similar to Fig. 2. It indicates that noisy entangled states can also behave as powerful as maximally entangled state in a sharing scenario. This observation can be important from the perspective of experiments where currently maximally entangled states can only be prepared with a certain but high visibility. Dark lines correspond to the scenario when Bob performs Bell-basis measurements while gray lines represent the elegant joint measurements in Eq. (6) by Bob. Both the axes are dimensionless.

Charus, the correlators take the form as

$$\begin{aligned}
 \langle A_x \rangle &= \langle B_y \rangle = \langle C_z^n \rangle = \langle A_x C_z^n \rangle = 0, \\
 \langle A_x B_y \rangle &= -\frac{v_1}{2} \cos \theta \delta_{x,y}, \\
 \langle B^y C_z^n \rangle &= \frac{v_2}{2} G'_n \cos \theta \delta_{y,z} \prod_{i=1}^{n-1} K'_i, \quad \text{where } K'_i = \frac{1 + 2F'_i}{3} \\
 \langle A_x B^y C_z^n \rangle &= -\frac{v_1 v_2}{2} G'_n (1 + \sin \theta) \prod_{i=1}^{n-1} K'_i \text{ if } xyz \in \{123, 231, 312\} \\
 &= -\frac{v_1 v_2}{2} G'_n (1 - \sin \theta) \prod_{i=1}^{n-1} K'_i \text{ if } xyz \in \{132, 321, 213\} \\
 &= 0 \quad \text{otherwise.}
 \end{aligned} \tag{37}$$

Finally, the corresponding bilocal expression reads

$$\begin{aligned}
 \mathcal{B}_{\mathcal{E}}(A^1, B, C^n) &= \frac{\cos \theta}{2} \left[v_1 + v_2 G'_n \prod_{i=1}^{n-1} K'_i \right] \\
 &\quad + 3v_1 v_2 G'_n \prod_{i=1}^{n-1} K'_i.
 \end{aligned} \tag{38}$$

In the unidirectional case, when $v_1 = v_2 = v$, to show the violation of bilocal models at round n , the weak measurement

parameter has to satisfy

$$G'_n > \frac{6 - v \cos \theta}{6v^2 + v \cos \theta} \prod_{i=1}^{n-1} \frac{1}{K'_i}. \quad (39)$$

If we consider $\theta = 0$, one can show that the measurement settings of Alice and Charus considered above is optimal [44].

Maximally entangled state as resource. When $v = 1$, i.e., the resource state is maximally entangled, at most two Charus can violate (7) to share bi-nonlocality with a single Alice. Here, using $G_1^{'cr} = \frac{5}{7} \approx 0.714$, we find that $G_2^{'cr} \approx 0.893$.

Werner state as resource. As shown in Fig. 3, when $EoF > 0.935$, two Charus can sequentially violate bilocal inequality involving EJM with a single Alice, while even a single Charu cannot violate bilocal model for $EoF < 0.689$.

If the shared state is NME, we find that the situation is much more involved. Taking $\theta = 0$ (i.e., each elements in the basis contains a very small amount of entanglement, $E = 0.355$), we observe that two Charus can violate the corresponding bilocal model with a single Alice sequentially only when $E_{in} > 0.998$. Notice that a single Alice-Charu duo cannot show bi-nonlocality when $E_{in} < 0.976$. Notice, moreover, that even for a singlet states as initials, very less number of Charus can exhibit bi-nonlocality with A^1 in comparison with BSM reported in the preceding section. It can be argued that such a disadvantageous situation emerges since each element of $EJM|_{\theta=0}$ contains a very low entanglement value, $E = 0.355$ compared to the elements of BSM, having unit entanglement. It seems that to obtain the violation of $\mathcal{B}_{\mathcal{E}}$, there is a competition between the entanglement content of the shared states and joint measurement basis and the choice of the optimal measurement strategies by Alice-Charu pair.

VI. DISCUSSION

In recent times, it has been established that unsharp measurements can provide certain benefits in quantum information processing tasks which cannot be reached by using projective (sharp) measurements due to its trade-off nature between the disturbance on the system and information obtained from the system. One prominent example is the sharing of entangled states in time-like separated observers which can be confirmed via the violation of Bell inequality, entanglement witnesses, steering inequality, etc. The violation of bilocal models of the resulting states after observers perform unsharp measurement sequentially are applied to detect nonlocality in the sharing scenario.

Two kinds of sharing scenarios are considered – unidirectional protocol where one of the observers performs unsharp measurement, and bidirectional process in which both the observers perform unsharp measurements. In the unidirectional scenario, we found that a maximum of six observers can exhibit bi-nonlocality when the shared state is maximally entangled. The maximum number of rounds for which the sharing of entangled states can be detected via the violation of bilocal models decreases with the decrease of entanglement content

of the initial shared states. We also observed that there exists a critical entanglement value of entanglement above which the multiple rounds of sharing bi-nonlocal states are possible. The situation changes drastically in the bidirectional case. In particular, the maximum number reduces to two when the shared state is close to maximally entangled states. After completion of our work, we notice that when both the observers share (noisy) maximally entangled states and want to employ network nonlocality by performing unsharp measurements in star and chain networks, a maximum of two rounds of detection for network nonlocal correlations is reported [43] (cf. also [42]).

We also showed that the number of rounds where the sharing is possible also depends on the measurement performed by the middle party in an entanglement swapping experiment. Specifically, we found that instead of Bell-basis measurement, the elegant joint measurement [44] has destructive effects on the protocol. It is possibly due to the fact that the elements of the elegant joint measurement except Bell-basis measurement contains a low amount of entanglement. We demonstrated that for a fixed elegant joint measurement, the maximum of two rounds of sequential sharing is possible even in the unidirectional situation when the shared state is either maximally entangled or close to the maximally entangled (in terms of non-maximally entangled pure state and maximally entangled state admixed with white noise).

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Appendix A: Discussion: weak measurement

The notion of weak measurement can be captured by the formalism of unsharp measurement [33] in two outcome measurement scenario using the set of positive operator values measurement (POVM) or effective operators denoted as $E_{\lambda}^a = (\mathbb{I} + (-1)^a \lambda n^i \sigma^i)/2$ with $i = 1, 2, 3$, $a \in 0, 1$ and $\lambda \in (0, 1]$. Each POVM element can be written as a sharp projector mixed with white noise as

$$\begin{aligned} E_{\lambda, \vec{n}}^a &= \lambda P_{\vec{n}}^a + \frac{1 - \lambda}{2} \mathbb{I}, \\ P_{\vec{n}}^a &= \frac{\mathbb{I} + (-1)^a n^i \sigma^i}{2}, \\ E_{\lambda, \vec{n}}^0 + E_{\lambda, \vec{n}}^1 &= \mathbb{I}. \end{aligned} \quad (A1)$$

In this formalism, the outcome independent unnormalized state of the system after measurement according to the Luder transformation rule, can be written as

$$\begin{aligned}\rho' &= \mathcal{W}_{\vec{n}}(\rho) = \sqrt{E_{\lambda, \vec{n}}^0} \rho \sqrt{E_{\lambda, \vec{n}}^0} + \sqrt{E_{\lambda, \vec{n}}^1} \rho \sqrt{E_{\lambda, \vec{n}}^1} \\ &= \sqrt{1 - \lambda^2} \rho + \left(1 - \sqrt{1 - \lambda^2}\right) (P_{\vec{n}}^0 \rho P_{\vec{n}}^0 + P_{\vec{n}}^1 \rho P_{\vec{n}}^1).\end{aligned}\quad (\text{A2})$$

Here, we identify $\lambda = G$ as the precision of the measurement and $\sqrt{1 - \lambda^2} = F$ as the quality factor or the disturbance generated on the state due to the performance of the measurement. The optimal pointer condition for information gain-disturbance trade-off is automatically satisfied by the unsharp formalism given by $F^2 + G^2 = 1$.

In the same way, we can get the outcome dependent unnor-

malized post measurement state as

$$\begin{aligned}\rho' &= \mathcal{W}_{\vec{n}}^a(\rho) = \sqrt{E_{\lambda}^a} \rho \sqrt{E_{\lambda}^a} \\ &= \frac{F}{2} \rho + \frac{(1 + (-1)^a G - F)}{2} (P_{\vec{n}}^0 \rho P_{\vec{n}}^0) \\ &\quad + \frac{(1 - (-1)^a G - F)}{2} (P_{\vec{n}}^1 \rho P_{\vec{n}}^1).\end{aligned}\quad (\text{A3})$$

1. Recursion relation of BRGP function in unidirectional case

For simplicity, the general BRGP expression can be expressed for the rounds, $n = 1, 2, 3$ which finally leads to the recursion relation, given in Eq. (23) and the condition for sequential sharing. They are given by

$$\begin{aligned}\mathcal{B}(A^1, B, C^1) &= \sqrt{G'_1} \left\{ \sqrt{|\cos \theta_1 \cos \phi|} + \sqrt{|\sin \theta_1 \sin \phi|} \right\}, \\ \mathcal{B}(A^1, B, C^2) &= \sqrt{G'_2} \left\{ \sqrt{|\cos \theta_2 \cos \phi((1 + F'_1) + (1 - F'_1) \cos 2\theta_1)|} + \sqrt{|\sin \theta_2 \sin \phi((1 + F'_1) - (1 - F'_1) \cos 2\theta_1)|} \right\}, \\ \mathcal{B}(A^1, B, C^3) &= \\ &\sqrt{G'_3} \left\{ \sqrt{|\cos \theta_3 \cos \phi((1 + F'_1)(1 + F'_2) + (1 - F'_1)(1 + F'_2) \cos 2\theta_1 + (1 + F'_1)(1 - F'_2) \cos 2\theta_2 + (1 - F'_1)(1 - F'_2) \cos 2\theta_1 \cos 2\theta_2)|} \right. \\ &\left. + \sqrt{|\sin \theta_3 \sin \phi((1 + F'_1)(1 + F'_2) - (1 - F'_1)(1 + F'_2) \cos 2\theta_1 - (1 + F'_1)(1 - F'_2) \cos 2\theta_2 + (1 - F'_1)(1 - F'_2) \cos 2\theta_1 \cos 2\theta_2)|} \right\}.\end{aligned}\quad (\text{A4})$$

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